# Finding Approximate Shape Regularities for Reverse Engineering

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# Abstract

Current reverse engineering systems can generate boundary representation (B-rep) models from 3D range data. Such models suffer from inaccuracies caused by noise in the input data and algorithms. The quality of reverse engineered geometric models can be improved by finding candidate shape regularities in such a model, and constraining the model to meet a suitable subset of them, in a post-processing step called beautification. This paper discusses algorithms to detect such approximate regularities in terms of similarities between feature objects describing properties of faces, edges and vertices, and small groups of these elements in a B-rep model with only planar, spherical, cylindrical, conical and toroidal faces. For each group of similar feature objects they also seek special feature objects which may represent the group, e.g. an integer value which approximates the radius of similar cylinders. Experiments show that the regularities found by the algorithms include the desired regularities as well as spurious regularities, which can be limited by an appropriate choice of tolerances.

**Keywords:** Beautification; Shape Regularities; Similarity; Reverse Engineering; Geometric Interrogations and Reasoning.

## 1 Introduction

Reverse engineering shape has a variety of applications in design and manufacturing, such as redesign. We aim to create a system that reconstructs a boundary representation (B-rep) model from a simple engineering object with a minimum of human interaction. It should be suitable for naive users and non-specialists as well as engineers. The generated model should have all the intentional shape regularities present in the original design.

A valid B-rep model can be generated from dense 3D range data obtained from multiple views of an object using a laser scanner [1, 2]. The multiple views are merged into a single 3D point set, which is triangulated and segmented into subsets representing the faces of the object [3]. To each subset, a surface approximating the points is fitted. These surfaces are stitched to form a valid B-rep model. For this project we only consider planar, spherical, cylindrical, conical and toroidal surfaces which either intersect at sharp edges or are connected by fixed-radius rolling ball blends. Reliable surface fitting methods exist for these surfaces [4] and many interesting engineering objects can be generated using only such surfaces [5]. For our current project we focus on objects with up to about 200 primary faces, which is a realistic limit achievable with current reverse engineering technology.

Sensing errors from the data acquisition and numerical errors arising from the successive algorithmic steps distort the created model. We must also consider additional errors introduced by possible wear of the object and the particular manufacturing method used to make the object. Even with increased accuracy of sensing techniques and surface fitting methods, some errors may still remain.

Our approach to creating an ideal model from an approximate initial model uses geometric reasoning to detect approximate shape regularities and impose them on the model. One way of doing this is to use constrained surface fitting methods [6, 7]. This might, for instance, require that two planes are fitted simultaneously under the constraint that they are orthogonal. A second approach is to identify features like slots and pockets whose approximate location and type are provided by a human being and use this information to improve the results of the segmentation and surface fitting phase [8].

Our approach tries to improve the approximate models without human assistance in a post-processing step called *beautification*. Instead of improving the model during surface fitting, we analyse the generated B-rep model to find approximate regularities and adjust the model accordingly (see Fig. 1). Given sufficient constraints generated from the regularities, an ideal model may be deduced from the constraints without further use of the measured point data.

This paper presents methods to find approximate shape regularities in B-rep models. Later work will address the constraint imposition strategy. Our shape regularities are defined in terms of similarities between properties of B-rep model elements, and similarities between these properties and special properties. For instance, we look for approximately equal radii of cylinders. We also try to find a special value, e.g. an integer, which is approximately equal to the average radius. The regularities are local in the sense that they relate to single properties of one or a small number of B-rep model elements. A different approach takes a global view of looking for symmetries of the complete model [9].

We assume that the desired regularities in the initial model are sufficiently distinct from the noise in the model. Our methods produce a large set of possible regularities the ideal model might possess instead of a small set of very likely regularities. These potential regularities will probably contain mutually inconsistent constraints, so a subsequent decision process must choose a maximal, consistent set of regularities to enforce on the model.

The following section introduces the types of regularity detected by our methods. Then we discuss the notion of similarity as the fundamental concept used by our methods. In the remaining sections we describe our algorithms for finding regularities, and show how they perform using several examples.

#### 2 Shape Regularities and Feature Objects

We express shape regularities in terms of properties of face, edge and vertex groups of the B-rep model. A property of a group of one or more of these elements is represented by a typed *feature object*. The *type* of a feature object depends on the property it represents. Feature objects are stored as vectors of some dimension d depend-

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Shana Daramatara	Equal shape perspectors	5		
Shape I arameters	S. Equal stable parameters.			
	Special values for shape parameters.			
	Simple integer relations between shape parameters.	4		
Axis Directions	Parallel axis directions.	5		
	Sets of axis directions with the same angle relative to a special direction (axis directions on planes and cones).	4		
	Symmetrical arrangements of axis directions.	4		
Axes	Axes intersecting in a point.	3		
	Aligned axes.	3		
	Parallel axes arranged along lines and grids with regular distances between them.	3		
	Parallel axes arranged symmetrically on cylinders.	2		
Positions	Equal positions.	2		
	Regular distances between positions arranged on a line, a 2D grid or a 3D grid.	3		
	Equal positions under projection.	3		

Table 1 Shape Regularities



Fig. 1 Reverse Engineering Phases

ing on the type. For example, the axis of a cone or cylinder forms a feature object of type *axis* represented by a direction and a position. The radius of a cylinder and the semi-angle of a cone are two other types of feature object. A single B-rep model element can generate several different feature objects of various types. Note that further feature objects can be created from groups of feature objects, or from groups containing feature objects and B-rep model elements. For instance, axis feature objects may generate intersection points as further feature objects.

The regularities we seek are derived from a survey [5] identifying those commonly present in simple engineering parts. An overview is given in Table 1. The number in the last column indicates how common the regularity is, with 5 meaning very common to 1 meaning rare.

We seek approximately equal shape parameters from surfaces and edges, such as radii, lengths and angles. Further we try to find special values for these parameters, like integers and integer multiples of  $\pi$ , and simple integer relations of the form  $n_2p_1 = n_1p_2$ between shape parameters  $p_l$  with integers  $n_k$ .

Directions in a B-rep model such as surface normals and directions of axes (referred to as axis directions) are clustered into parallel sets. In addition we look for special directions with which all elements of a set of axis directions make the same angle. Sets generated in this way are axis directions lying in a plane or on a cone. Furthermore they may be arranged symmetrically (see Fig. 2).

For axes, i.e. axis directions associated with a position, we seek common intersection points and aligned axes. For parallel axes, we seek regular arrangements of these axes along a line or on a grid with regular spacing. Parallel axes may also be arranged symmetrically on a cylinder.

We check if positions obtained from vertices, centres of spheres and tori, apices of cones, axis intersection points, etc. have approximately equal locations or are regularly arranged on a 2D or 3D grid. We also seek positions which are equal after projection onto a special plane or line derived from the main axis directions present in the model.

#### **3** Similarity and Special Feature Objects

We define similarity measures indicating how close two feature objects are to each other. They are used by our hierarchical clustering algorithm to find similar feature objects and to find special feature objects similar to a given feature object. For a set X of feature objects of the same type, a similarity measure is a symmetric, nonnegative function  $\delta : X \times X \to \mathbb{R}^+_0$ , such that  $x_1 = x_2$  implies  $\delta(x_1, x_2) = 0$ . We call two feature objects  $x_1, x_2 \in X \varepsilon$ -similar (with respect to  $\delta$ ) if  $\delta(x_1, x_2) < \varepsilon$ , ( $\varepsilon \in \mathbb{R}^+$ ).

This definition suffices as a measure to decide if a feature object is close to a special feature object. However, for the clustering algorithm  $\delta$  should be a similarity metric, i.e. it should also fulfil the triangle inequality and  $\delta(x_1, x_2) = 0$  should imply that  $x_1 = x_2$ . Appropriate similarity measures and metrics are defined later for the different regularities.

To represent groups of similar feature objects by a single feature object, we need an averaging method avg to merge similar feature objects of the same type. Given two  $\varepsilon$ -similar feature objects  $x_1$  and  $x_2$  of the same type and two positive weights  $\omega_1, \omega_2$ , we generate a new average feature object  $x_{avg} = avg(x_1, \omega_1, x_2, \omega_2)$ , which represents  $x_1, x_2$  such that  $\delta(x_{avg}, x_l) < \varepsilon$ , (l = 1, 2).

We use a hierarchical clustering algorithm to find similar feature objects of the same type. Given a set of feature objects  $X = \{x_l\}$ , a similarity metric  $\delta$ , and an averaging method avg, we create a partition  $\{C_k\}$  of X such that each cluster  $C_k$  contains feature objects which are  $t_a$ -similar for some tolerance  $t_a$ .  $C_k$  is represented by some feature object  $c_k$  and a tolerance  $t_k$  as the maximum distance between  $c_k$  and the elements of  $C_k$ . Note that we do not limit the number of clusters  $C_k$ , but generate as many clusters as required by the tolerance  $t_a$ , which limits the width of the clusters.

To form a hierarchical structure, we create a nested subset structure for the clusters  $C_k$ . We partition a  $C_k$  further into disjoint

Element	(a) Shape Parameter	(b) Type	(c) Axis Direction
plane	—	—	normal
sphere	radius	length	—
cylinder	radius	length	direction of the axis
cone	semi-angle	angle	direction of the axis
torus	major radius	length	direction of the axis
	minor radius	length	
straight edge	distance between end points	length	direction of the line between the end-points
circular edge	radius	length	normal of plane in which the circle lies
	angle of the circle segment	angle	
elliptical edge			normal of plane in which the ellipse lies

Table 2 Shape Parameters and Axis Directions

sub-clusters  $S_j$  represented by feature objects  $s_j$  and with a tolerance  $t_j$  such that  $\delta(s_j, x_l) < t_j$  for  $x_l \in S_j$ . The sub-clusters must be sufficiently distinct from each other, i.e. for each pair  $S_{j_1}$ ,  $S_{j_2}$ ,  $(j_1 \neq j_2)$ , the condition  $\delta(s_{j_1}, s_{j_2}) - t_{j_1} - t_{j_2} \ge t_d$  must be fulfilled for some tolerance  $t_d$  satisfying  $0 < t_d \le t_a$ . The sub-clusters are again partitioned recursively if possible.

The hierarchical clustering reveals groups of feature objects in a cluster which are considerably closer to each other than to the other elements of the cluster. This can be used in the subsequent steps to decide which regularities should be enforced and makes the tolerances maximal rather than optimal. Maximal tolerances ensure that all intended regularities are found, but they also cause the detection of regularities which are not present in the ideal model and which could be avoided by an optimal choice of tolerances.

Various approaches to the clustering problem exist [10]. The most common is the agglomerative technique, which starts with the smallest value of  $\delta(c_l, c_k)$ , and combines the two elements to form a new element  $\hat{c}$  which replaces  $c_l$  and  $c_k$ . Clusters and sub-clusters are repeatedly formed until only one cluster remains, or the distance between the clusters is too large. A brute force solution searching for the closest elements each time requires  $O(n^3)$  time.

To improve this method, we use a matrix containing the distances between feature objects of the same type, and maintain a quad-tree like data structure to keep track of the closest pair [11]. The distances  $\delta(c_l, c_k)$  between the elements  $c_l$  and  $c_k$  representing clusters are stored in a matrix D for l < k. The clusters are grouped arbitrarily into pairs  $(C_l, C_{l+1})$  and the distance between two pairs is defined to be the minimum distance between the four clusters of the two pairs. These pairs define a new closest pair problem of half the size, which can be solved recursively. The closest pair in D is at the root of this recursive data structure. After initializing, each update operation needs to update at most two rows and two columns of each of the matrices. Thus, the clustering requires  $O(n^2)$  space and time assuming that  $\delta$  and avg require constant time and space.

Additional methods find special feature objects similar to an average feature object, e.g. the average radius of a group of similar cylinders could be approximately an integer. Note that there might be more than one such value within a given tolerance. In the following sections details of clustering and finding special values are given for different types of feature object.

#### 4 Shape Parameters

The first feature objects we consider are shape parameters (see Table 2(a)). These describe the shape of an element independently of its location and orientation. Each shape parameter is treated as a separate feature object and more than one shape parameter feature object may arise from a single element.

To find similar parameters we use our hierarchical clustering algorithm. It is not very useful to compare parameters like the semi-



Fig. 2 Planar and Conical Angle-Regular Directions

angle of a cone with the radius of a sphere, so we assign a type to each parameter (see Table 2(b)) and only cluster parameters of the same type. The values for angle and length parameters are in different units, so the tolerances used depend on the parameter type. We use  $t_{la}$  and  $t_{ld}$  for lengths and  $t_{aa}$  and  $t_{ad}$  for angles. The averaging method computes the weighted average of real numbers independent of the parameter type.

We assume that fixed-radius rolling ball blends are identified in the B-rep model, and thus we handle them as edge attributes and not as surfaces. We treat blend radii as separate shape parameters of type *blend* and find similar and special values for blend radii.

Note that the parameters from edges can depend on surface parameters. For example, the radius of a circle which is the boundary of the top of a cylinder has to be equal to the cylinder radius. Such relationships are better handled in the constraint solver phase where they should become obvious.

Clustering shape parameters results in a hierarchical structure of similar shape parameters sorted by type. In the following subsections we present methods to find special parameter values and special integer relations between parameter values. Both problems reduce to finding a simple fraction approximating a real number.

**4.1 Special Parameter Values and Integer Relations.** For each average shape parameter for a cluster or sub-cluster we try to find a simple fraction approximating its value. As there might be more than one special value within a given tolerance, we create a list of appropriate special values for each shape parameter.

Let v be an average shape parameter from a cluster or subcluster. Depending on the parameter type we also have a tolerance  $t_a$  (either  $t_{la}$  or  $t_{aa}$ ) and we choose a family of functions  $f_l : \mathbb{R} \to \mathbb{R}$ which represents the scales on which we look for simple fractions. For angle parameters we use the two functions  $f_{\pi} : v \mapsto v\pi$  and  $f_t : v \mapsto \arctan(v)$ . For length parameters the family is defined by  $f_{K_l} : v \mapsto vK_l$  where the  $K_l$  are base units for length measurements like 1.0, 0.1 or 2.54 (cm to inch conversion).

To find special values for v, we look for integer pairs  $n_k$ ,  $m_k$ such that  $v \approx f_l(n_k/m_k)$  for all specified functions  $f_l$ .  $s_k = f_l(n_k/m_k)$  is a valid approximation of v if  $|s_k - v| < t_a$ . Hence, to find a special value for v, we try to find fractions  $n_k/m_k$  which approximate  $w_l = f_l^{-1}(v)$  for each function  $f_l$  within a tolerance

- I. The function has been called as  $rec_frac(x, n, m, neg_frac, M)$  where x is the value which has to be approximated by a fraction, n, m are the two integers representing the fraction n/m found so far, neg\_frac indicates if x has to be added or subtracted from n/m and M is the maximum denominator allowed at this recursion level with the initial limit being  $M_0$ .
- II. Let b = 1.
- III. While the denominator a = round (b/x) is not larger than M:
  - 1. If a > b:
    - A. Let r = x b/a.
    - B. If r is negative, set neg\_r to true and r = -r. Otherwise set neg\_r to false.
    - C. If neg\_frac is true, then the new numerator is p = na mb, otherwise p = na + mb.
    - D. The new denominator is q = ma.
    - E. If r < t, then add p/q to the list of special values, if it is not already in it.
    - F. If  $r > t_{rmin}$  and p/q was not already in the list of special values, call rec\_frac(r, p, q, neg\_r,  $M_0M$ ). 2. Let b = b + 1.

## Algorithm 1 Recursive Algorithm rec\_frac for Finding Simple Fractions

Table 3 Finding Special Values for 0.63 with t = 0.05,  $M_0 = 5$  and  $t_{\text{rmin}} = 0.01$ 

t derived from  $t_a$ . For linear  $f_l : v \mapsto vK_l$  we use the tolerance  $t = t_a/K_l$ . For  $f_t$ , we use  $t_{aa}$  as tolerance with the condition  $||v| - \pi/2| > t_{aa}$  assuming that  $v \in [-\pi, \pi]$ . The algorithm for determining simple fractions for  $w_l$  with tolerance t is described in Section 4.2. Note that when considering multiple functions  $f_l$ , we have to compare the results for different functions and eliminate duplicates, e.g.  $\arctan(1) = \pi/4$ .

The algorithm to find simple fractions approximating a real number can also be used to find integer relations between shape parameters. For each pair of clusters  $C_k$  and  $C_j$  represented by  $c_k$  and  $c_j$ , we look for a set of integer relations  $n_l/m_l$  which approximate  $w = c_k/c_j$ . To have consistent input for the algorithm, we order the clusters such that  $c_k > c_j$ . If  $|n_l/m_l - w| < t_{ra}$  for some tolerance  $t_{ra}$ , we accept  $n_l/m_l$  as a valid approximation for w. To avoid finding equality relations, i.e. 1/1 relations, we only generate relations between clusters and sub-clusters which do not belong to the same cluster.

**4.2 Finding Simple Fractions.** The previous subsection reduced finding special parameter values and special integer relations between parameters to finding a list of special fractions for a real number w in the open interval (w - t, w + t) with tolerance t. We assume that integer values are always special for w and that w is non-negative.

Although finding integer relations between real numbers [12] and recognizing numerical constants [13] are well studied related problems, such methods assume that the real numbers are exact or that a very close approximation is required. We try to find special fractions which lie in an interval containing w.

If 1/(2t) < m, then more than one n/m for a fixed m could be in the interval (w - t, w + t). To avoid this ambiguity we set the maximum allowed value for m to  $M_0 = \text{floor}(1/(2t))$ . To limit the number of possible special fractions found,  $M_0$  should not be larger than  $M_{\text{max}}$ , say 10. Also note that if t is larger than 1/2, we would ignore the integers, so  $M_0$  should at least be 1. Optionally, if we have to work with large t, we can set a minimum  $M_{\text{min}}$  for  $M_0$ which is larger than 1.

If  $M_0$  is small, we could simply multiply w by each  $m \le M_0$ and check if |w - n/m| < t with n = round(wm), i.e. check if the relation n/m is within the tolerance limits. However, for large  $M_0$ , this becomes expensive as we must check many fractions for a large number of real numbers w.

To overcome this problem, we expand the simple method with a continued fractions approach [14] as listed in Algorithm 1. We start by approximating w by  $a_0 + x_0$  where  $a_0 = \text{round}(w)$ . The error  $x_0$  is approximated by using the simple method recursively. Starting with  $b_l = 1$  at recursion level l, we get  $x_l = b_l/a_l + x_{l+1}$  for  $a_l = \text{round}(b_l/x_l)$  and error  $x_{l+1}$ . We get additional approximations by increasing  $b_l$  by one, until  $a_l$  is larger than some limit M. The process is repeated recursively for each error  $x_{l+1}$ until the error is smaller than some tolerance  $t_{\text{rmin}}$ . Table 3 contains an example for 0.63 with t = 0.05,  $M_0 = 5$  and  $t_{\text{rmin}} = 0.01$ where the special values generated are marked with  $[\rightarrow \cdot]$ . For the intermediate result 1/2 we do one recursion step to find special values for the error 0.13.

To use this method for an arbitrary real number, we must do some preprocessing. First its sign and the closest integer  $a_0$  to w have to be found. If  $x_0 = |w - a_0| < t_{\rm rmin}$ , the integer is the only special value for w and we stop immediately. Otherwise we continue with  $x_0$ . To process small denominators, we seek simple fractions for  $1 - x_0$  if  $x_0 < 0.5$ , as functions of the type  $b/x_0 - \operatorname{round}(b/x_0)$ used in the recursive algorithm are more likely to generate small denominators for  $x_0 \in (0.5, 1)$ . Based on the tolerance t for w, we also generate an initial limit  $M_0$  for the allowed denominator as explained earlier.

To find simple fractions for  $x_0$ , we call the recursive algorithm rec\_frac listed in Algorithm 1, where we must take into account the preprocessing steps to find the final special values. In recursion step l we solve the problem to find simple fractions p/q for a real number  $x_l$  with q < M using the simple method, where the limit M for the denominator depends on the recursion level starting with  $M = M_0$ . For  $x_0$  we already have a simple fraction n/m such that  $x_0 = n/m \pm x_l$ . Whether  $x_l$  has to be added or subtracted from n/m is indicated by the flag neg\_frac. For each  $b = 1, 2, \ldots$  we check the denominator for  $x_l$ . Depending on neg\_frac we add or subtract it to n/m, generating a new approximation p/q for  $x_0$ . If this approximation is close enough, it is added to the list

of special values. If the error r = |x - b/a| is still larger than some tolerance  $t_{\rm rmin}$ , and p/q was not already in the list of special values, we call rec\_frac for r recursively with a new limit  $M_0M$ for the denominator. By multiplying M with the initial limit  $M_0$ , we ensure that we can still find fractions for a smaller real number within the denominator limit. We always reduce fractions n/m to simplest terms, in order to keep the values small and to ensure that we add *simple* fractions to the list of special values.

Note that the algorithm can only miss special values close to w if their denominator is larger than  $M_0$ . The precision increases with the depth of the recursion.

#### 5 Axis Directions

Directions arising from the B-rep model, like plane normals, provide the basis for another class of regularities. Note that some directions are also associated with a position. In this section we are only interested in the angular arrangement of the directions. Directions with positions are covered in Section 6.

We extract unit vectors representing direction feature objects from B-rep model elements (see Table 2(c)). Opposite directions, i.e. the unit vectors d and -d, are identified. This space of directions is the real projective plane  $P_2$  and can be represented by the unit sphere with antipodal points identified. We call its elements *axis directions*. Regular arrangements of the axis directions correspond to points and circles in  $P_2$ . The way in which the directions are arranged on the circles might create further regularities.

By using the hierarchical clustering algorithm, we can find parallel axis directions represented by points in  $P_2$ . As a similarity measure, we use the smaller angle between axis directions  $d_1$  and  $d_2$ ,  $\angle(d_1, d_2) = \arccos(|d_1^t d_2|)$ . The weighted average between two axis directions  $d_1$  and  $d_2$  with weights  $\omega_l$  is  $(\omega_1 d_1 + \operatorname{sign}(d_1^t d_2)\omega_2 d_2) / (\omega_1 + \omega_2)$ . The two tolerance values required for clustering are  $t_{aa}$  and  $t_{ad}$  defined earlier. The resulting parallel axis direction clusters represent the main directions present in the model. Even if the number of directions is large, we expect to find only a limited number of *different* directions.

The following subsections discuss the arrangements represented by axis directions lying on circles in  $P_2$  and regular arrangements on these circles. Axis directions that are on a great circle of the sphere lie in a plane. Axis directions that are on a small circle of the sphere represent axes that are on a cone. The arrangement of the axis directions in the plane or the cone can be symmetric, which we call planar or conical angle-regular (see Fig. 2).

**5.1 Regular Arrangements of Axis Directions.** A set of axis directions  $\{d_l\}$  on a circle satisfies the equation system  $|d_l^t x| = a$ , where  $x \in P_2$  is the centre of the circle. For a = 0, the directions lie in a plane with normal x, which we call an axis plane. For  $a \in (0, 1)$ , we have a cone with axis x, which we call an axis cone. Note that for a plane, taking the absolute value of  $d_l^t x$  is not required, and we can drop it for a cone if all axis directions have the same orientation relative to x.

To find the sets of axis directions  $d_l$  lying in an axis plane, we cluster the normals representing all axis planes generated from each pair of linearly independent axis directions. The clustering is done in the same way as clustering parallel axis directions, but we employ the clustered parallel axis directions instead of all axis directions. This not only reduces the number of normals generated, but also avoids approximately linearly dependent axis directions. We also only consider the main clusters of parallel axis directions and not the sub-clusters.

A similar method is used to find axis cones. For each triple  $d_1$ ,  $d_2$ ,  $d_3$  of axis directions representing parallel axis direction clusters, we generate an axis cone with axis direction c and semi-angle  $\alpha$  by solving the linear equation system  $d_l^t x = 1$ , (l = 1, 2, 3), from which we get  $\alpha = \arccos(|x^t d_1|/||x||)$  and  $c = \alpha x$ . We have

to avoid flat axis cones that actually represent axis planes or axis cones that represent parallel axis directions by rejecting axis cones for which  $\alpha < t_{\rm aa}$  or  $|\pi/2 - \alpha| < t_{\rm aa}$ , as the  $d_l$  are not exact.

The axis cones represented by pairs  $(c_l, \alpha_l)$  are clustered using  $\sqrt{\angle (c_1, c_2)^2 + (\alpha_1 - \alpha_2)^2}$  as similarity measure and the tolerances  $t_{aa}$  and  $t_{ad}$ . The averaging method for two cones generates an average cone from the weighted average of the axis directions and the weighted average of the semi-angles.

Note that we cluster the description of the axis planes or cones, and not the axis directions used to create them. For instance, three axis directions  $d_1$ ,  $d_2$  and  $d_3$  generate three axis planes  $p(d_1, d_2)$ ,  $p(d_1, d_3)$  and  $p(d_2, d_3)$ . If  $p(d_1, d_2)$  and  $p(d_1, d_3)$  are combined into a single cluster,  $p(d_2, d_3)$  is not automatically added to the same cluster. As  $p(d_1, d_2)$  and  $p(d_1, d_3)$  are only approximately the same plane,  $p(d_2, d_3)$  is not necessarily the same plane. Hence, the way the axis planes are generated creates dependencies between them and the clusters might not have a transitive structure with respect to the generating axis directions. We get similar dependencies for the axis cones.

Any reasonable choice of tolerances should ensure that dependent axis planes or cones are in the same main cluster even if they are in a different sub-cluster. If this is not so, then the tolerances are too restrictive for the errors in the model or they are too relaxed, causing the combination of axis directions which are not related.

Handling these dependencies in the clustering algorithm is rather expensive as it requires adding all other clusters dependent on the new cluster when combining two clusters to form a new one, and combining the clusters could no longer be done in constant time. Instead, the dependencies are handled by a post-processing step which combines dependent (sub-)clusters.

We can further improve the approximation representing the regular arrangements once we know the sets of axis directions forming them. Given a set of axis directions  $d_l$  which are approximately on a circle in  $P_2$ , we get a linear equation system  $d_l{}^t c = \alpha$  for l = 1, 2, ..., n. For  $\alpha \neq 0$ , i.e. for axis cones and parallel axes, we can solve the above system in a least-squares sense [15]. Using this approach for axis planes creates a constrained optimization problem  $||[d_1{}^t c; ...; d_n{}^t c]||_2 \rightarrow \min$  for c subject to the condition  $||c||_2 = 1$  (or at least  $||c||_2 \gg 0$ ) which is more difficult to solve.

**5.2** Angle-Regular Axis Directions. Examining axis planes and cones further can reveal symmetrical arrangements of the directions (see Fig. 2). Given a set of axis directions in a plane or on a cone, we look for subsets such that the angles between the axis directions are integer multiples of a base angle  $\beta$ . The subsets can be incomplete, i.e. not all multiples of  $\beta$  have to be present. The problem is to identify appropriate subsets of multiples which approximately satisfy this condition. We first describe the problem for axis planes, then generalise it for axis cones and other regularities introduced later.

Let  $\{d_j\}$  be a set of m axis directions in an axis plane and let  $\alpha_{lk}$  be the angle between  $d_l$  and  $d_k$ . We call the  $d_j$  angle-regular if there is a  $\beta \in \{\alpha_{lk}\}$  such that  $\beta = \pi/n$  for  $n \in \mathbb{N}$  and for each  $\alpha_{lk}$  there is an integer p such that  $\alpha_{lk} = p\beta$ . This means that the angle between some reference direction  $d_{j_0} \in \{d_j\}$  and each other direction is an integer multiple of  $\pi/n$ . As we identify opposite directions, we only consider angles in the interval  $(0, \pi]$ . Based on which multiples of  $\beta$  are present, we decide whether an angle-regular set is considered to be a regularity. We avoid too small base angles  $\beta$  by setting a maximum  $N_{\text{max}}$  for n, which should be smaller than  $\pi/(2t_{\text{aa}})$ . Note that in the approximate case the base angle  $\beta = \pi/n$  is only close to some  $\alpha_{lk}$ .

At this stage we do not look for base angles which are not approximately present as an angle between axis directions. For instance, if we have two approximate angles  $2/6 \pi$  and  $5/6 \pi$  relative to some direction, the underlying base angle  $\pi/6$  is not detected.

- I. Let  $\{d_i\}$  be a set of m objects for which distance-regular subsets with respect to the condition reg<sub>n</sub> are sought.
- II. Compute the distances  $\alpha_{lk}$  between objects  $d_l$  and  $d_k$  for l < k < m.
- III. For all reference objects  $d_{j_0}$  with  $j_0 < m$  and for all distances  $\alpha_{lj_0}$  with  $j_0 < l < m$ :
  - 1. Find the candidates  $\beta_{j_0n}$  for some  $n \in \mathbb{N}$  such that  $\operatorname{reg}_n(\beta_{j_0n}, \alpha_{l_j})$  is true.
  - 2. Add  $\beta_{j_0n}$  to the list of candidates for  $d_{j_0}$  unless it is an integer multiple of one of the  $\beta_{j_0n_0}$  already in the list.
  - 3. If  $\beta_{j_0n}$  has been added to the list, remove any  $\beta_{j_0n_0}$  from the list which are integer multiples of  $\beta_{j_0n}$ .

IV. For each reference object  $d_{j_0}$  with  $j_0 < m - 1$  consider the subset  $\{d_{j_0}, \ldots, d_m\}$  and for each candidate  $\beta_{j_0 n}$  with  $n = 1, \ldots, N_{j_0}$ : 1. For each object  $d_j$  with  $j_0 < j \le m$ :

- A. If for  $f = \alpha_{lj_0} / \beta_{j_0 n}$ , we have  $| \operatorname{round} (f) \beta_{j_0 n} \alpha_{lj_0} | < t_a$ , then do:
  - a. Record object  $d_j$  to be the multiple round (f) of the base distance. It is stored as a multiple round (f) under the following conditions.
  - b. If there is already an object o as the multiple round (f) and o is the parent of  $d_j$ , then only replace o by  $d_j$  if  $d_j$  is sufficiently closer to the multiple as indicated by  $t_{pd}$ .
  - c. If there is already an object o as the multiple round (f) and  $d_j$  is the parent of o, then only replace o by  $d_j$  if o is not sufficiently closer to the multiple as indicated by  $t_{pd}$ .
  - d. If there is already an object o as the multiple round (f) and  $d_j$  is not related to it in the hierarchical structure, then replace o by  $d_j$  if  $d_j$  is closer to the multiple.
  - e. If there is no object o as the multiple round (f), then make  $d_j$  this object.
- 2. If the arrangement in the distance-regular subset for  $\beta_{j_0n}$  is regular (see text), note it as a regularity. In addition remove integer multiples of  $\beta_{j_0n}$  from the base distance candidate list for the reference objects  $d_l$  which are in the current regular subset.

#### Algorithm 2 Finding Distance Regularities

PSfrag replacements



Fig. 3 Planar Angle-Regular

An efficient implementation for these cases is left as future work. One might, for instance, use approximate integer relations.

From an abstract point of view, the  $\{d_j\}$  are objects which can be ordered with respect to an arbitrary reference object  $d_{j_0}$  such that we get a set of distances  $\alpha_{lj_0}$  between  $d_{j_0}$  and  $d_l$  for all l. To find a base distance  $\beta$  similar to the base angle in the approximate case, we look for all  $\beta$  fulfilling a regularity condition  $reg_n(\beta, \alpha)$ for an  $\alpha \in \{\alpha_{lk}\}$ . This condition depends on the object type. In this more general case we call the set distance-regular.

Given a set of *m* objects  $\{d_j\}$  we seek a minimum number of subsets which are approximately distance-regular with respect to reg<sub>n</sub> (see Algorithm 2). First we compute all distances  $\alpha_{lk}$ , (k < l < m), between the objects, where a column  $j_0$  in that matrix with elements below the diagonal represents the distances to the reference object  $d_{j_0}$ . From these distances we derive a set of possible  $\beta_{j_0n}$  for each  $d_{j_0}$  using the condition reg<sub>n</sub>. Because we seek approximate distance-regular arrangements, a single  $\alpha_{lk}$  can generate more than one base angle candidate depending on the tolerances. Finally we try to find distance-regular subsets by checking the distances  $\alpha_{lj_0}$  for each reference object  $d_{j_0}$  and all base distance candidates  $\beta_{j_0n}$ .

The distances between the objects used in the first step depends on the object type. For axis directions in a plane,  $\alpha_{lk}$  are the angles between the axis directions  $d_l$  and  $d_k$ . We must take care to choose the angle that lies consistently to the right of the reference object  $d_l$  with respect to the normal q of the axis plane (see Fig. 3). With  $v = q \times d_k$  we have  $\alpha_{lk} = \arccos(d_k^{t}d_l)$  if  $v^t d_l \ge 0$  and  $\alpha_{lk} = \pi - \arccos(d_k^{t}d_l)$  if  $v^t d_l < 0$ . This also allows us to identify which of the  $k\pi/n$ , (k = 0, ..., n - 1), directions a particular  $d_j$ occupies for some  $n \in \mathbb{N}$ .

In the next step, candidates  $\beta_{j_0n}$  for base distances are derived from each  $\alpha_{lj_0}$  such that  $reg_n(\beta_{j_0n}, \alpha_{lj_0})$  is satisfied for some  $n \in \mathbb{N}$ . A particular  $\beta_{j_0n}$  is added to the list of base angle candidates for  $d_{j_0}$  unless it is an integer multiple of some other  $\beta_{j_0n_0}$  already in the list. If any  $\beta_{j_0n_0}$  in the list is an integer multiple of  $\beta_{j_0n}$ , we remove  $\beta_{j_0n_0}$ . This ensures that only the smallest  $\beta_{j_0n}$  is used to produce a minimum number of distance-regular subsets. For the axis planes, we use the regularity condition  $|\beta_{j_0n} - \alpha_{lk}| < t_{aa} \operatorname{as reg}_n(\beta_{j_0n}, \alpha)$  to find all  $\beta_{j_0n} = \pi/n$  such that  $n < N_{\max}$ .

In the final step we check each set  $\{d_{j_0}, \ldots, d_m\}$  for  $j_0 \ge m-1$  for distance-regular subsets with respect to the candidates  $\beta_{j_0n}$ . The reference object is always an element of a distance-regular subset. Therefore, for each  $\beta_{j_0n}$  we search for objects in  $\{d_{j_0+1}, \ldots, d_m\}$  that form an approximately distance-regular subset. A  $d_l$  for  $l \in \{j_0 + 1, \ldots, m\}$  belongs to the distance-regular subset if  $| \operatorname{round} (f_l) \beta_{j_0n} - \alpha_{lj_0} | < t_a$  for  $f_l = \alpha_{lj_0} / \beta_{j_0n}$ . However, we only allow one object for each multiple of  $\beta_{j_0n}$ , as the objects were already clustered. Thus, if we have two objects  $d_{l_1}$  and  $d_{l_2}$  for the same multiple, we use the one closer to p, unless one of the objects is a parent of the other in the hierarchical clustering structure. Then we take the parent, even if the child is closer, if the difference between the two tolerances is smaller than some  $t_{pd}$ .

Before we accept a distance-regular subset of objects as a regularity we consider which multiples of  $\beta_{j_0n}$  are present. We accept a planar angle-regular subset as a regularity if either all multiples, or at least every second one is present, or at least three consecutive multiples are occupied. If a set is accepted, we must remove multiples of  $\beta_{j_0n}$  in the list of candidate base distances for all objects  $d_l$  in the distance-regular subset, otherwise subsets of the distanceregular subset found will be identified as distance-regular later.

For axis planes we check in addition if the angle between two axis directions, which is not involved in any angle-regular subset, has a special value (see Section 4.1).

To find angle-regular axis directions on an axis cone, we again use Algorithm 2. The definition for a conical angle-regular subset is similar to the planar case. We project the axis directions from the cone onto the plane through the origin orthogonal to the axis of the cone. After projection, opposite directions on the plane represent different axis directions on the cone, and so we must use base angles of the form  $2\pi/n$ . However, each axis direction on the cone can still point in one of two directions. Thus, we always project the direction pointing to the same side of the plane defined by cone normal c, i.e.  $d_l^{t}c > 0$ . Note that three axis directions forming an orthogonal system generate a special conical angle regularity with a base angle  $\pi/2$  on a cone with semi-angle  $\arccos(1/\sqrt{3})$ .

#### 6 Axes

Axis directions which are also associated with a position represent axes. For these, we consider the arrangement of the positions as well as the direction information. We seek approximate intersections of axes and regular arrangements of parallel axes.

For cylinders, cones and tori we use the root point of the surface as the position. The root point of a cone is its apex and the root point of a torus is its centre. The root point of a cylinder is an arbitrary point on the central axis. For elliptical edges we choose the centre of the ellipse, and for straight edges we choose an arbitrary point on the edge as the associated position.

Planar faces do not have an obvious root point, but one can be defined by considering the boundary loops. Reasonable choices for root points are the average of the vertex positions for each loop and the centre of the convex hull of each loop. Note that these define multiple axes for a planar face. Other possibilities exist for defining root points of planar surfaces and more general types of edges.

First we search for intersection points of axes. The approximate intersection point of each axis pair is found by computing the minimum distance between them. If it is smaller than  $t_{la}$ , we say the axes intersect and use the mid-point of the shortest line between them as the approximate intersection point. These intersection points are clustered to find sets of axes intersecting approximately in a point, where we use actual axes and not axis direction clusters. We only intersect axes belonging to different axis direction clusters to avoid trying to intersect approximately parallel axes.

Furthermore, we seek regular arrangements of parallel axes. For each cluster of parallel axis directions we extract the elements having a root point. In addition we consider the axes from conical angle-regular arrangements if an intersection position for such axes has been found. To explore regular arrangements of the parallel axes, we project the locations for the axes onto a plane through the origin orthogonal to the axis direction.

Finding regular arrangements of axes is now reduced to finding regular arrangements of points in a plane. By clustering the points using the tolerances  $t_{la}$  and  $t_{ld}$  we detect approximately equal axes. The resulting clusters are examined further to detect if the points lie on a line, a grid or a circle, and are regularly spaced.

**6.1 Regularly Arranged Parallel Axes.** We detect regular arrangements on lines and grids by generating a line for each pair of points, clustering these lines into approximately parallel lines and finding approximately equal lines in the parallel line clusters. The equal line clusters are then checked for regular point arrangements. By examining pairs of approximately orthogonal groups of parallel lines, we find the grids.

First we generate a line for each pair of points  $p_1$ ,  $p_2$  in the plane, represented by the vector  $d = p_2 - p_1$ . To cluster the lines into parallel groups, given two such vectors  $d_1$ ,  $d_2$  with  $||d_2|| > ||d_1||$ , we use the similarity measure  $\delta(d_1, d_2) = ||d_1 - (d_1^t u)u||$  with  $u = d_2/||d_2||$ . By using the distance between  $d_1$  and its projection onto  $d_2$  instead of the angle between the two vectors, we also take the distance between the points into account. If we used the angle between the vectors alone, two points which are on different parallel lines in a grid and sufficiently far apart might generate a line which is approximately parallel to the grid lines. Similarly, we also define the weighted average between the directions,

$$avg_{p1}(d_{1}, \omega_{1}, d_{2}, \omega_{2}) = \frac{\|d_{1}\|\omega_{1} + \|d_{2}\|\omega_{2}}{(\omega_{1} + \omega_{2})^{2}} \\ \left(\frac{d_{1}}{\|d_{1}\|}\omega_{1} + \operatorname{sign}(d_{1}{}^{t}d_{2})\frac{d_{2}}{\|d_{2}\|}\omega_{2}\right).$$
(1)

Each of the resulting clusters of approximately parallel lines is again clustered to get groups of equal lines. To do this, the lines in a cluster of parallel lines are represented by the two points  $p_1$ ,  $p_2$ . The tolerance used is  $t_{la}$ , and the similarity measure between two lines represented by  $p_1$ ,  $p_2$  and  $q_1$ ,  $q_2$  is based on the average distance between two corresponding end-points. If  $(q_1 - q_2)^t (p_1 - p_2) \ge 0$ , the points  $p_1$  and  $q_1$ , and  $p_2$  and  $q_2$  correspond to each other. Otherwise  $p_1$  and  $q_2$ , and  $p_2$  and  $q_1$  correspond. Let  $d_1$  and  $d_2$  be the two distance vectors between corresponding points and d be the unit vector representing the direction of the parallel line cluster. Then the similarity measure between the two lines is  $(||d_1 - (d_1^t d)d|| + ||d_2 - (d_2^t d)d||)/2$ , i.e. the average of the orthogonal distance with respect to d between the corresponding points. The weighted average of these two lines is the line with end-points  $r_l = p_l + \omega_2/\omega_1(d_l - (d_l^t d)d), (l = 1, 2)$ .

For each group of approximately equal lines we seek base distances such that the distances between a subset of the points on the line can be represented as integer multiples of a base distance, analogously to the angle-regular arrangements using Algorithm 2. The main difference is that base distances are not regular values such as  $\pi/n$ , but any distance can be a base distance. This simplifies the algorithm as we get at most one possible base distance per distance between the points. Note that we split approximately equal lines into different lines if we find different base values such that each line represents one base value.

After these steps we have sets of parallel lines, where a line might be marked as distance-regular. To generate grids we search for orthogonal pairs of distance-regular parallel line sets. Lines which are not distance-regular or for which we cannot find an orthogonal partner are noted as simple regularities. For the orthogonal pairs of line sets we try to find regular grids.

Each orthogonal pair is handled separately. First the two sets of parallel lines in the pair are grouped into lines with compatible distance-regular arrangements. Two parallel lines belong to the same distance-regular group if one of the base distances is approximately a multiple of the other, and the distance between the two reference points on the line in the parallel direction is approximately an integer multiple of the smaller base distance. This produces two lists of groups containing compatible distance-regular lines. Corresponding elements of each group generate grids in such a way that the distances between the lines in one group fit on the distanceregular arrangement of the other group and vice versa.

The generated grids are not unique because various diagonals of a grid can form a distance-regular line and combining orthogonal pairs of these diagonals can form additional grids. We must thus find and remove these diagonal lines and grids and add additional points on them to the fundamental grid. A line is a diagonal of a grid if the reference point of the line lies on the grid and the base distance  $d_l$  of the line can be generated from the two base distances  $d_1$  and  $d_2$  of a grid. Let  $D_j$  be the unit vector representing the directions for the distance  $d_j$  in the grid and  $D_l$  be the direction of the line. The line is a diagonal of the grid, if

$$d_l D_l \approx \operatorname{round}\left(\frac{d_l D_l^{\,t} D_1}{d_1}\right) d_1 D_1 +$$

$$\operatorname{round}\left(\frac{d_l D_l^{\,t} D_2}{d_2}\right) d_2 D_2.$$
(2)

Another grid with base distances  $d_3$ ,  $d_4$  and corresponding directions  $D_3$ ,  $D_4$  is a diagonal of the grid if its reference point is on the grid and the distances are compatible. Without loss of generality, we assume that  $d_1^2 + d_2^2 < d_3^2 + d_4^2$ , i.e. the diagonal of the second grid is longer than the diagonal of the first one. Then the second grid is a diagonal if Eq. (2) holds for  $d_l = d_3$ ,  $D_l = D_3$  and  $d_l = d_4$ ,  $D_l = D_4$ . The grid with the shorter diagonal is the fundamental grid and we eliminate the one with the longer diagonal.



Fig. 4 Example Objects

Any distance-regular lines not on a grid and not removed as diagonals of a grid are noted as regularities. In addition we also check whether the base distances have a special value for all distanceregular lines and grids (see Section 4.1).

Parallel axes arranged along a circle can be found by generating all circles formed by triples of points in the plane and clustering the circles.

# 7 Positions

This section briefly discusses regularities of positions such as vertices, root points of surfaces, and axis intersection points. Using the hierarchical clustering algorithm for this points with tolerances  $t_{la}$  and  $t_{ld}$ , the Euclidean distance as similarity measure, and the weighted average between points, we get approximately equal positions.

The resulting position clusters could be examined further for distance-regular arrangements on lines and 2D and 3D grids using a similar approach to the previous section. However, as the points are in  $\mathbb{R}^3$ , many additional options are present, and a general search for such arrangements is expensive. Furthermore, finding unintended approximately regular arrangements is likely to happen often. Thus, we limit the search to special lines and grids. If a B-rep model has one or more orthogonal systems, we search for distance-regular arrangements on lines and grids build from the directions of the orthogonal system(s). If the model has a main axis without an orthogonal system, we search for distance-regular arrangements on lines parallel to the main axis and on 2D grids orthogonal to it.

These directions are used further to find partially equal positions, i.e. positions which are equal under projection. We cluster positions projected on a plane through the origin orthogonal to a special axis direction. Adding a second orthogonal axis direction, we project the positions further onto a line to find partially equal positions with respect to two axis directions.

# 8 Examples

The methods described in this paper were implemented<sup>1</sup> on a GNU/Linux platform (with a 450Mhz Pentium III and 256MB of RAM), and tested using various simple objects and more complex objects from [16]. Errors induced by the reverse engineering process were simulated by translating and rotating faces of the exact models before generating a point cloud. Models reconstructed from point clouds from both exact models, and models perturbed

by varying amounts, were analysed. A test was considered to be successful if all the intended regularities in the model were found when using tolerances agreeing with the amount of perturbation.

Some examples are shown in Fig. 4. Execution times for models with up to 200 faces range from a few seconds up to a few minutes. **R** gives the total number of regularities found by the methods. **U** is the number of unwanted regularities detected, i.e. those which are not part of the design and conflict with the desired regularities. **M** is the number of important regularities missed by our algorithms as identified by a human being. These values were determined for objects with angles perturbed by  $3^{\circ}$  and positions perturbed by 0.3 units, with angle tolerances of  $5^{\circ}$  and length tolerances of 0.5 units; the number of unwanted regularities can be reduced by more careful tolerance value selection. It appears that all tolerances for the methods can be expressed in terms of angular and positional tolerances which could be derived for a particular reverse engineering system from the errors in the reconstruction of a known test object. We now briefly describe the main results for the example objects.

The planar angle-regular axis directions with base angle  $\pi/4$  of object (a) were detected. Axes of the slots at the top and the bottom, and the axes of the planar faces were found to intersect; the intersection points are partially equal. Various aligned axes, and two orthogonal systems, were also found as conical angle regularities. Object (b) consists of a five-sided truncated pyramid with a six-sided truncated pyramid cut out of it. Both conical angle regular arrangements were found. Some of the planar faces of the two pyramids were considered to be parallel depending on the perturbation and the tolerance values. In the perturbed model additional conical angle-regular arrangements were found involving either the top of a pyramid and its sides or a combination of the faces of the two pyramids. Those caused by the perturbation were easily identifiable as they contradicted the two main conical angle regularities.

The main axis for each of objects (c) and (d) was detected as a set of parallel axis directions and aligned axes. For object (c), the axes of the holes arranged on a circle and the axes of the holes intersecting on the main axis were detected as well as some other axis intersection points. Some cylinder radii were unwantedly considered to be equal. In object (d) two planar angle regularities of the faces orthogonal to the main axis with base angle  $\pi/2$  and the angle  $\pi/6$  between the two arrangements were detected. A number of vertices and intersection points of axes close to each other were unwantedly considered to be equal. Some regular arrangements of axes of planar faces were not detected or incorrectly determined; some parameters were considered to be equal, causing some special parameter values not to be found. Object (e) was found to have a single orthogonal system; multiple distance-regular arrangements

<sup>&</sup>lt;sup>1</sup>The sources are available at <uri: http://www.langbein. org/software/sil/>.

of axes along lines were also found. The type of regularities which caused problems were very similar to those for object (d). Nevertheless, nearly all expected regularities were correctly found.

Preliminary tests of partial models actually reverse engineered from point data gave similar results to the simulated data. Suitable tolerance values for real data appear to depend mainly on the reverse engineering system and not on the particular object.

Generally, all regularities found were either present in the model or could be understood by considering the perturbation and the tolerance values used. With small tolerance values, only a small set of very accurate and thus very likely regularities are found. With bigger tolerance values, at some point we get all the desired regularities. However, the number of unwanted regularities also increases, so that when all desired regularities are found, we also get some unwanted regularities. An automated decision about which regularities should be used to improve the model has to be made by a subsequent step. For the analyser the user has to supply the tolerances, which should be maximal rather than optimal due to the hierarchical structure of the regularities.

# 9 Conclusions

We have presented algorithms to find local shape regularities in terms of similarities between feature objects derived from B-rep model elements. Tests with various perturbed mechanical components were promising in the sense that the main approximate regularities of the parts were found quickly while few unwanted regularities were reported. Thus, the analysis methods provide a strong basis for our subsequent beautification steps.

In future work we will develop a constraint solving strategy for finding and imposing a maximal, consistent set of constraints containing the main regularities of a part to generate an ideal model.

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